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In re the Application of:

Kathy Maida-Smith and Steven W. Engle

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Title:

NETWORK SECURITY DATA MANAGEMENT SYSTEM AND METHOD

Box Patent Application Assistant Commissioner for Patents Washington, D.C. 20231

Sir:

NEW APPLICATION TRANSMITTAL

Applicants are enclosing an executed provisional application for filing in the U.S. Patent and Trademark Office (PTO), including twenty-four (24) pages of specification and abstract, and nine (9) pages of informal drawings.

This application is being filed without the filing fee under 37 C.F.R. § 1.53. Accordingly, applicants are believed to be entitled to a filing date based upon applicant's deposit of the attached specification, which is being filed under Certificate of Mailing § 1.10, with Express Mail #EL455184579US. Therefore, it is respectfully requested that a filing date of November 22, 1999, be granted this application. The filing fee will be timely filed.

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STATEMENT CLAIMING SMA (37 CFR 1.9(f) & 1.27(c)) S	ALL ENTITY STATUS MALL BUSINESS CONCERN	Docket Number (Optional) 66997-0102					
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PATENT APPLICATION

of

Kathy Maida-Smith Steven W. Engle

Entitled:

"NETWORK SECURITY DATA MANAGEMENT SYSTEM AND METHOD

CERTIFICATE OF MAILING 37 C.F.R. § 1.10

I hereby certify that this correspondence is being deposited with the United States Postal Service with sufficient postage as Express Mail, Express Mail No. EL455184579US, addressed to: BOX PATENT APPLICATION, Assistant Commissioner of Patents, Washington, D.C. 20231; on November 22, 1999.

Michael Hawes

TECHNICAL FIELD OF THE INVENTION

The present invention relates in general to monitoring data networks and more particularly to a data management system and method for compiling and displaying network security data.

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Network providers want the advantages of accessing outside resources, such as those available on the Internet, but do not want those contacts to result in threats of unauthorized information release, modification of internal records, or network downtime. They also need to protect the network from unauthorized actions performed by internal users. In order to counter those problems, a typical information technology network may include many network components that collect security data or perform functions safeguarding the confidentiality, integrity, or availability of the network, its attached systems, application software, and data. Examples of such network components include firewalls, proxy servers, intrusion detection systems, routers, and availability monitors. Each of those network components either collects or has access to data that is useful to network security and administration personnel.

Collecting and using the security data available in a typical network may be difficult and time consuming. The data provided by each network component may be organized into a series of categories that are inconsistent with the categories used by another network component. Even when identical categories can be identified, the data may be stored in different formats. For example, one number might be represented in floating point format while another number corresponding to the same quantity may be represented in fixed point format.

As a result of the dissimilar organization and formats of the security data and the resulting time and effort that would be required to transform it into a usable form, the network components that make up a typical network may not be configured to compile and store this data. Later, if a confidentiality, integrity or availability problem is suspected of having occurred, this data would not be available to confirm this. If the data was compiled and stored, it can be accessed in an attempt to reconstruct the relevant time period. The data from each network component is typically analyzed separately, though even the data from a single network component may be difficult to analyze because of an upgrade or change in software that changed the data output format. Thus, network security personnel typically do not have the resources to monitor the contemporary security data available

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from their networks. Even if a security concern is noted, the historical data may be available only on a component-by-component basis, if at all.

From the foregoing, it may be appreciated that a need has arisen for a system that compiles security data available from network components. A need has also arisen for a system that parses differently-organized data into records having a consistent labeling and access structure. Further, a need has arisen for a system that manages data having several formats. Additionally, a need has arisen for a system that stores the information used to parse and format different data so that it can be updated, if necessary, upon a software change or upgrade. A need has also arisen for a system that can display contemporary security data from the network components in response to database queries. Each of those needs is independent and can be addressed without addressing the other identified needs.

In accordance with the present invention, a method and system for managing network security data is provided, embodiments of which substantially eliminate or reduce at least some of the disadvantages and problems associated with typical systems for handling security data.

According to one embodiment of the present invention, a method for building a system including a database of data templates is provided that includes identifying sets of data categories, each set corresponding to security data received from one of the network components. Database record definitions are constructed with subdivisions matching the sets of data categories. Parser scripts are written that receive security data from network components and output records corresponding to the database record definitions. The parser scripts are then stored.

According to another embodiment of the present invention, a system for compiling security data from an information network is provided. That system includes at least two network components, each providing data in a different format. A data parser is coupled to the network components. The data parser has access to two parser scripts and is operable to produce categorized data from the network components' data using the parser scripts.

According to another embodiment of the present invention, a method for compiling network security data is provided that includes collecting the data from at least two network components. Parsing scripts corresponding to the network components are accessed and applied to the data from the network components to produce categorized and formatted data which is then stored.

Embodiments of the method and system for managing network security information of the present invention provide various technical advantages over typical security information systems and methods. For example, one technical advantage is facilitating real time access to security information. Another technical advantage is in reducing the need for extensive security personnel, each monitoring the output of different security network components. Yet another technical advantage is in allowing automated detection of events defined by information from multiple network components. Still another technical advantage is allowing relational database queries of security information from multiple network components. Another technical advantage is facilitating updates of information

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defining the data structures used by network components. Another technical advantage is facilitating the compilation of security data from a network whose network components are often replaced. Some of the embodiments of the invention may not provide every technical advantage identified. No one technical advantage is an essential element of the invention. Other technical advantages are readily apparent to one skilled in the art from the following figures, description, and claims.

BRIEF DESCRIPTION OF THE DRAWINGS

For a more complete understanding of the present invention and the advantages of the embodiment thereof, reference is now made to the following description taken in conjunction with the accompanying drawings, wherein like reference numerals represent like elements, in which:

FIGURE 1 illustrates a block diagram of an information technology network that produces security data and includes a system for compiling that security data;

FIGURE 2 illustrates a flow chart of a method for building and utilizing a security data database;

FIGURE 3 illustrates a block diagram of a system for compiling and displaying security data from an information network;

FIGURE 4 illustrates a data flow chart for a firewall component with an information technology agent;

FIGURE 5 illustrates a data flow chart for an intrusion detection system with an information technology agent;

FIGURES 6 and 7 illustrate a meta-data table update prepared in accordance with one embodiment of the invention.;

FIGURE 8 illustrates a data relationship chart of an uninstanciated meta-data database, an instanciated meta-data database, and global database tables;

FIGURE 9 illustrates transaction data received from a network component; and FIGURE 10 illustrates the data of FIGURE 9 in a global database table.

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DETAILED DESCRIPTION OF THE INVENTION

Figure 1 is a block diagram of an information technology network 10 that produces security information. Information technology network 10 includes network rings 12 that connect network components. Those network rings 12 can be of any of the many types known to network engineers, including token rings and fiber hubs. The network rings 12 can include connections to user workstations 14 and file servers 16. Intranet services are can be made available to user workstations 14 by file transfer protocol (FTP) servers 18 and web servers 20. Routers 22 analyze and direct information packets propagating between network components or to the Internet 24. Firewalls 30 can be included to control access into the network 10 and to control outside use of an Internetavailable FTP server 26 and web server 28. A proxy server 32 can be provided in order to monitor internal requests for information from the Internet 24 and can be configured to block some of those requests based on various criteria. Each network ring 12 can also be coupled to security network components that perform intrusion detection 34, availability monitoring 36, utilization monitoring 38, and alerting 40. An availability monitoring component 36 can monitor the availability of the network, of a host, of an application, or of a combination of those. Utilization monitoring 38 can also apply to any or all of those three targets. The security of an information technology network encompasses the confidentiality, integrity, and availability of the software and data that it and its attached network components store, transfer, and process. Some network security components require or can be used with separate management stations 42. In an embodiment of the present invention a system for compiling security data 44 is coupled to the information technology network 10 such that security data can be received from the network components over the network 10.

Figure 2 is a flow chart of an embodiment of the invention illustrating a method for building and utilizing a security data database. Three of the possible methods by which network component output information can be obtained are demonstrated. First, a specification for the product in question can be obtained. From that specification a person with information technology network experience can identify the categories of data, that the network component outputs, for example in a syslog or via Simple Network Management Protocol (SNMP). Individual elements within the

network component's output can be located based on absolute position within the line or a record

A meta-database is a database containing data describing other data. In one embodiment that database stores information which indicates the categories of output data from a network component and the format of data in those categories, providing a template of the output. The idea of a template can be explained in reference to a letter template on a word processor that stores information about where on the page the address and date are located, as well as the format in which the date and address are presented. For example, a date can be either in European format, where the day precedes the month, or American format, where the month precedes the day. The actual database software used to store the meta-data can be Oracle, SQL Server, or other well-known databases.

A second way of obtaining network component output information utilizes the Simple Network Management Protocol (SNMP) and a Management Information Base (MIB). A MIB, provided by the manufacturer of the network component in the form of a flat file, describes the data that can be extracted from the network component via SNMP and documents the syntax for extraction. Software MIB compilers are available for converting a MIB into a software application's internal form. This software application would then be capable, using SNMP, to extract data from the corresponding network component. A compiler which converts the MIB into a table update file appropriate for the meta-database can be programmed and is called an MIB integrator. Applying the MIB integrator to the MIB is called integrating the MIB and results in a table update including attribute data, syntax, and identification of the type of network component to which it corresponds. That table update can be input directly into the meta-database.

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A third method of obtaining attribute data can be used where the data is provided from the network component by an information technology agent (ITA). See the discussion of Figures 4 and 5 below for a more specific discussion of ITAs. In that case the output of the ITA must be analyzed to determine the categories and formats in which the data is transmitted. An ITA can also be documented, such that the documentation will define the categories and formats. Some network components will transmit security data through both their native systems and an ITA. In such a case, the product specification and the ITA output characteristics must be used.

Through use of the three methods already identified, a meta-database can be built and maintained. The attribute data in the meta-database is noninstanciated. Noninstanciated data is data that does not correspond to output from specific, physically present network components and instead corresponds to output from types of network components. For example, the attribute data for all routers having a certain product number is noninstanciated. The identical attribute data that corresponds to a router having that product number but also having a specific serial number is instanciated. The advantage of maintaining a meta-database of noninstanciated attribute data will become apparent. One embodiment of employing the attribute data to compile security data requires that the data be instanciated. Other embodiments that access the noninstanciated meta-database for compiling security information can also be employed.

One way to instanciate the attribute data comprises compiling a list of network components with instance data. Instance data identifies which specific network components from which security data is desired are present in the information network 10. For example, the list might include routers 22 having a certain product number and identify two serial numbers of routers of that type that are in the network and from which security information is desired. Once a network component list with instance data is obtained, the attribute data can be instanciated.

Instantiating the data includes producing a group of attribute data records, which includes at least records corresponding to network components from which security data is desired that are physically present in the information network 10. The records are produced by comparing a network component listed in the device list to the network components for which the noninstanciated meta-

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database has attribute data, identifying the attribute data corresponding to that network component, and creating a new record for each instance of that network component. That new record would include both the attribute data and the instance identification, one possibility of identification is a serial number.

For example, if the device list includes a firewall with a specific product number and corresponding instance data showing two such firewalls from which security data is desired currently coupled to the network, then two new records of instanciated attribute data would be produced. The firewalls would share a system type id but would have different system ids. A system table and interface table would be built for each physical firewall. Figure 8 further details the tables built for instanciated meta-data. In on embodiment, the instanciated data for each firewall would reference the attribute data from the noninstanciated meta-database through the system type id, which is used to identify both instanciated and uninstanciated meta-data. Each would also include an identifier such as a system id.

The advantage of maintaining the noninstanciated meta-database can now be seen. If the processes of compiling a device list with instance data and instanciating the attribute data are both automated, changes in networks components from which security data is desired can be implemented automatically on a routine basis. If a proxy server 32 is switched for another model, a new instanciated meta-database can be created automatically as long as the uninstanciated meta-database includes attribute data for that new type of proxy server.

The instanciated meta-database is maintained and provides information for two functions. First, together with uninstanciated meta-data, it provides the information necessary to choose the parameters of the global database tables in which the security information will be stored. Second, it determines what parser scripts are necessary for data received from the network components and what uninstanciated meta-data must be accessed to construct those parsers.

In one embodiment, a global database table is built for each transaction of a network component from which security data is desired. For example, a specific type of firewall can provide two different security data outputs. If the network includes two physical firewalls of that type, the

global database can include four tables. The tables are uniquely labeled in accordance with the system id and xaction id (transaction id). The instanciated meta-data associates a system type id with each physical network component from with data is desired. The uninstanciated meta-data can, based on the system type id, provide the template for each data output produced by that network component. The information from the instanciated and uninstanciated meta-data that is necessary to build a global database table is shown in Figure 8.

Once the global database tables have been built, parser scripts can be produced for each output of a network component for which there is attribute data in the instanciated database. The flat file table and fixed form tables of Figure 7, from the uninstanciated meta-data, contain parameters from a which a parser script can be prepared. The parser scripts and their supporting constructs enable the system's input format independence. Each network component chosen to provide security data can transmit that data to the system for processing by the corresponding parser scripts. That transmission can be performed under several known protocols such as Simple Network Management Protocol (SNMP) or syslog. Generally the transmission protocols used by the network 10 to send other data are used for security data as well. In general, parser scripts for a particular network component reformat the data received from that network component so that it can be entered in the global database tables. If the output data from a network component does not include a category that is included in a global database table field, the parser script will reject the output. This filters out "garbage" data that does not need to be collected, processed, or archived. If the received data is not rejected, the parser script arranges the data to match the categories and formats of records in the corresponding global database table. Categorized data is data that conforms to a record in a database table. Other embodiments of parser scripts are possible, for example, a parser scripts could recategorize but not reformat data received from network components.

Once the parser, using the parser scripts, has manipulated the information received from the network components, it loads the parsed data into the global database tables. The data can immediately be added to the corresponding table because it has been parsed. The data in the global tables is consistently labeled and accessible.

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Before, or during, the global database load, the parser can transmit the parsed data to an event detector. The event detector compares the parsed data to one or more event definitions. The event definitions can be formulated independent of the structure of the data from the network components because the event detector receives parsed data. If the parsed data matches one of the event definitions, a signal is generated by the event detector indicating that the event occurred. The signal may be recorded, displayed to a network manager, or it may initiate an automatic change in the status of the network, among other possibilities. For example, an event might be an attempt to break into a computer system from a site on the Internet by exploiting a known computer system vulnerability. In that case connection and any further attempts to connect from that site could be blocked until the security vulnerability is resolved.

Figure 3 illustrates an embodiment of the invention in which a system compiles security information from network components through intermediaries and allows that information to be searched and displayed. The network components include a utilization monitor 38, firewall systems 30, intrusion detection systems 34, file servers 16, proxy server 46, and availability monitor 36. The file servers 16 can use various operating systems including UNIX and Windows NT. Information technology agents 46, ITAs, can be deployed at network components to collect and transmit information. Figures 4 and 5 illustrate operation of example ITAs. For example, if a network component does not have the capacity to output security information to which it has access, it may be beneficial to deploy an ITA 46 to get that information. The ITAs 46 and network components without ITAs can transmit data to distributed data managers 48, DDMs. DDMs 48 are most useful in very large networks containing a multitude of network components providing security data, but they can be used in any size of network. In one embodiment, each DDM 48 utilizes parser scripts for the network components from which it receives data. In another embodiment the DDM does not parse the data, but merely forwards it. In another embodiment, a DDM 48 is not used and all data is delivered directly to the Central Data Manager 50. DDMs 48 can forward data as they receive it or wait and forward data in batches according to the amount of data or a time interval. DDMs 48

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can compress data before forwarding it. DDMs 48 can also encrypt data before forwarding it. The DDMs 48 forward data to the Central Data Manager 50, CDM.

The data is received at the data interface 52. If the data is encrypted or compressed, the data interface 52 returns it to its original form. In an embodiment where the DDMs do not parse the data, the data interface 52 can be provided with parser scripts in order to parse the data. The data interface 52 can be coupled to a data storage 54 to which the parsed data is transmitted. Unparsed data can also be transmitted to the data storage 54. For more permanent storage the data, parsed and unparsed, can be transmitted to a tape library 56.

In one embodiment, data is delivered directly from the data interface 52 to a global database 58 for storage. The global database, when it is present, can include tables whose structures correspond to the categories and format of the parsed data as discussed with regard to Figure 2. An event detector 60 can receive parsed data either directly from the data interface 52 or from the global database 58. The event detector 60 compares the parsed data to one or more event definitions and generates an event signal upon a match.

Coupled to the global database is an database interface 62 for receiving and responding to queries. The queries can be real time, in which case the data corresponding to the query is provided in the form required by the query at webserver 64 for forwarding to a web browser 66 or the database can be configured to produce reports at specified times that are stored in the webserver 64. The webserver 64 can be accessed by a webbrowser 66 having a java interpreter 68. Software is available for handling queries from a webbrowser 64 for most major software databases.

Figure 4 illustrates one example of an information technology agent (ITA) 46 used with a firewall. Elements of the firewall, including firewall network packet filters 400 and firewall network application proxies 402, collect security information that is transmitted to the firewall audit log manager 404. Such information may include the time and data when a communication occurred, the source and destination ip addresses, the protocol used, the number of packets sent and received, and the number of bytes sent and received.

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A firewall may include system performance measurement utilities 408 and software to collect and log performance data 410 produced by the utilities. Such data may include the percentage of time that a CPU is idle, and the percentage of time it is used for each of system tasks, input and output tasks, and user tasks. Like the system log application 412, performance data may be provided directly to a DDM 48 or the CDM 50. Some performance data may be provided to the system log application 412.

The ITA 46 consolidates the data received from the firewall audit log manager 404, the system log application 412, and system performance measurement utilities 408. The ITA 46 may compress or encrypt 416 the data before sending it 418 to a DDM 48 or the CDM 50. In order to provide an accuracy check, the ITA 46 may also hash data and attach the hash result to the data. Hashing is running the file through a mathematical algorithm that yields a fixed-length value or key that represents the original file. The mathematical algorithm is the hashing function. A hashing function is secure if it is computationally infeasible to find a file that corresponds to a given value or key, or to find two different files which produce the same value or key. For example, the Secure Hashing Algorithm (SHA-1) was made available by the National Institute of Standards and Technology on April 17, 1995. The CDM 50 can, upon receiving the data, transform the file with the secure hashing function and check the key which was attached by the ITA 46. If the hashing function result matches the attached key, it is likely that the data was provided without corruption.

FIGURE 5 illustrates one example of an information technology agent (ITA) 46 used with an intrusion detection system (IDS). Elements of the IDS, including IDS Network Monitor 500 and IDS network event detector 502, collect security information that is transmitted to the IDS audit manager log 504. While an IDS may provide different types of information, the flow of data to the ITA 46 and DDM 48 or CDM 50 in the embodiment of Figure 5 is similar to the flow of data in Figure 4 for the firewall. Thus IDS elements 506-512 provide information that the ITA 46

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consolidates 514, encrypts 516, and transmits 518 to a DDM 48 or the CDM 50. The ITA may also compress or hash the data as discussed with reference to Figure 4.

Figures 6 and 7 illustrate one embodiment of an uninstanciated meta-database. A metadatabase includes several tables and updating a meta-database may involve adding records for each of those tables. The transaction table 700 includes records where each correspond to a record or group of data elements that is provided by one type of network component. The records include a 'name', identification code ('xaction id') and 'description' of the transaction. The records also identify the ids of the type of network component that provides the transaction, the 'system type id', and the source of the data associated with the transaction, 'stream id'.

The system type table 702 includes records that each correspond to a type of network component that provides at least one transaction. That descriptive information is also used in formulating the device list with instance data, see Figure 2, so that the meta-database can be instanciated. Once the descriptive information is used to identify the 'system type id', the transactions corresponding to that 'system type id' can be identified.

The stream table 704 includes records that specify the separation of streams by type, 'stream type'. For example, in a delimited flat file stream type elements are separated by a code. For that type of file, the stream identification table 704 includes that identification "DELIMITED FLAT FILE" and identifies the separator, "space" in the figure. A fixed form flat file uses designated, unchanging points in the file to separate elements. The records also include the operating system nomenclature for the data location or source identified as the stream path.

Moving to Figure 7, the element table 706 includes records for each element of each transaction. The records identify the transaction that includes the element, the 'element id', descriptive information regarding the element, and the format of the element. For example, the network transaction includes a date which uses the DATE_ISO format. Numerical data may also be additionally characterized as a count or a percentage. Each element record includes information about the updating of that element by the network component. Each element can be associated with

a stream through the transaction id, 'xaction id'. The transaction id is used to identify a record in the transaction table 700 which includes the corresponding 'stream id'.

The delimited flat file table 708 and fixed form flat file table 710 are both stream type tables and each includes records for the elements that are associated with a stream of that type. For example, because stream id T100204 is delimited flat file, the elements corresponding to that stream are listed in the delimited flat file table 708. Additional tables could be included for additional stream types.

The records in the stream type tables include the formatting information that identifies the element in the stream. That information is used to create the parsing scripts discussed in prior figures. For example, a parsing script for the fixed form flat file associated with xaction id X317189 is written to separate the first ten characters (each an ASCII 7-bit character) and insert them in the portion of the record labeled "Date" in the DATE_ISO format. The next eight characters are inserted in the "Time" portion of the record in TIME_24 format. The parser of transaction X200154, on the other hand, reads 7-bit ASCII until it finds as space. The characters prior to the space are then inserted as "Date".

Figure 8 illustrates a data relationship chart of an uninstanciated meta-data database, an instanciated meta-data database, and global database tables. The tables shown above the shaded bar are uninstanciated and correspond to the uninstanciated tables shown with records in Figures 6 and 7. The system table 712 and interface table 714 contain instanciated meta-data because they correspond to a network component that is physically present. Those tables are built from the device list with instance data discussed in Figure 2. The system table 712 contains a record for each network component in the network. The fields contain the attributes of the physical network components including system type id, system id, host name, domain name, description, default route, serial number, license key, organization, location, and contact. While actual network components may share a 'system type id', each has a unique 'system id'.

The interface table 714 allows characterization of network component interfaces. While a network component has only one record in the system table 712, it can have multiple records in the

interface table 714, each identifiable by the 'system id'. The interface table fields include 'system id', 'interface name', 'interface host name', 'interface IP address', 'interface network mask', 'interface side', 'interface type', and 'interface speed'.

Each 'system id' is combined with the corresponding 'xaction ids' to determine the names of the global database tables that must be built to receive data from the parsers for each system and transaction. Each global database table is labeled with the 'system id' and the 'xaction id' reflecting the transaction data for the system it will be storing. The global database table records include all the elements associated with the' xaction id' in the element table 706 formatted in accordance with the "legal value" and encoding portion of those records and labeled in accordance with the "name" portion of those records.

The global database table S0214_X200154 (718) corresponds to system id S0214, which corresponds to system type id T100001. It also corresponds to xaction id X200154 and the elements in its records are categorized to match the element records in the element table 706 that have X200154 in the xaction id field. As such it will be storing data for transaction 'X200154' for system 'S0214'. The second table 720 corresponds to the same physical network component, S0214, but to a different transaction, X317189, of that system type. Consequently, the records of the second table 720 have different fields because they correspond to element records that have X317189 in the xaction id field. Because both the parsers and the global database tables are constructed consistently with the meta-data, the output of the parsers can be directly added to the global database tables in real time.

Figure 9 illustrates transaction data received from a network component as a fixed form flat file and Figure 10 illustrates the data of Figure 9 in a global database table. The data parsers built in accordance with the present invention enable the consistent format of the global database tables. The data in the tables is labeled consistently regardless of the format in which the network component output the data. For example, the dates received in the X317189 transaction are labeled and formatted consistently with dates received in the X200154 transaction. As a result, the data can be subject to database table joins initiated by the relational database. The meta-data can be used to

structure the user interface and the database queries. The data can also be made available through

Thus, it is apparent that there has been provided, in accordance with the present invention, a network security data management system and method that satisfies the advantages set forth above. Although embodiments has been described in detail, it should be understood that various changes, substitutions, and alterations can be made herein without departing from the spirit and scope of the present invention as defined by the following claims.

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WHAT IS CLAIMED IS:

A method for compiling parser scripts each corresponding to the structure of security data received from a network component comprising the steps of:

- a) identifying sets of data categories, each set corresponding to security data received from one of a plurality of network components;
- b) constructing database record definitions, each defining a record subdivided in accordance with one of the sets of data categories;
- c) writing parser scripts that receive security data from the network components and output records, each record corresponding to one of the record definitions; and
 - d) storing said parser scripts.
- 2. The method of claim 1 further comprising the steps of:
 - e) determining the format of each category in said sets;
 - f) formatting the subdivisions to match the formats of the categories of the set to which the definition corresponds; and wherein

each of the output records of step (c) correspond in format to one of the record definitions.

- 3. The method of claim 1 further comprising the steps of:
 - e) building database tables in a relational database each having the fields of one of the database record definitions; and
 - f) inserting output records received from the parser scripts into the tables.
- 4. The method of claim 2 further comprising the steps of:
 - g) building database tables in a relational database each having the fields and formats of one of the database record definitions; and
 - h) inserting output records received from the parser scripts into the tables.

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- 5. The method of claim 1 wherein:
 - at least one of the sets of data categories is identified, at least in part, from the product specifications of the network components.
- 5 6. The method of claim 1 wherein:

at least one of the sets of data categories is identified, at least in part, by applying a Management Information Base (MIB) integrator to a Management Information Base for the corresponding network component.

An information network security data compilation system, comprising:

- a) a first network component;
- b) a second network component; and
- a data parser coupled to the first and second network components having access to a first parser script and a second parser script, the data parser is operable to produce categorized data from the data received from the first and second network components with the first and second parser scripts, respectively.
- 8. The data compilation system of claim 7 wherein:
 - a) the first network component is a firewall and
 - b) the second network component is an intrusion detection system.
- 9. The data compilation system of claim 7 further comprising:
 - a) a third network component and
 - b) a distributed data manager; and wherein:

the data parser is coupled to the second and third network components through the distributed data manager which collects and compresses data from the second and third network components and forwards the compressed data to the data parser.

- 10. The data compilation system of claim 7 further comprising:
 - a) a third network component;
 - b) a second data parser coupled to the third component having access to a third parser script, the second data parser operable to produce categorized data from the data received from the third network component with the third parser script; and
 - c) a relational database coupled to the first and second data parsers.
- 11. The data compilation system of claim 7 further comprising:
 - a) a display coupled to the data parser; and
 - b) a relational database coupled between the data parser and the display, and wherein: the data parser transfers the categorized data to the relational database.
- 12. The data compilation system of claim 11 wherein:
 the relational database receives a data query, and
 the display shows a portion of the categorized data, up to and including all the data, from the
 relational database, corresponding to the data query.
- 13. The data compilation system of claim 12 wherein:
 the data queries are submitted and the portions are shown through a web browser interface.
- 14. The data compilation system of claim 7 further comprising:
 - a) an event detector coupled to the data parser and wherein: the event detector compares the categorized data to a predetermined event definition and provides a signal if a match is found.

- 15. The data compilation system of claim 7 further comprising:
 - a) an information technology agent and wherein: the network component is programmed by software, the agent collects security data from the software, and the data provided from the first network component is the security data collected by the agent.
- 16. The data compilation system of claim 7 wherein: the data parser produces formatted and categorized data.
- 17. The data compilation system of claim 7 wherein:

 data from the first network component is security data and data from the second network component is security data.
- 18. The data compilation system of claim 7 wherein:

 data from the first network component is encrypted and decrypted.

A method of compiling network security data comprising the steps of:

- a) collecting security data from a plurality of network components;
- b) accessing a plurality of different parser scripts, each script corresponding to one of the network components;
- c) applying the plurality of different parser scripts to the security data to produce categorized and formatted data; and
- d) storing the categorized and formatted data.
- 25 20. The method of claim 19 wherein: the plurality of network components includes at least a firewall and an intrusion detection system.

- 21. The method of claim 19 further comprising the steps of:
 - e) transmitting the categorized and formatted data to a relational database;
 - f) providing a user interface for submitting queries to the relational database; and
 - g) displaying the categorized and formatted data, or a subset thereof, in accordance with submitted queries.
- 22. The method of claim 21 wherein:

step (e) occurs prior to step (d) and step (d) comprises storing the categorized and formatted data in the relational database.

- 23. The method of claim 19 further comprising the steps of:
 - e) comparing the categorized and formatted data to at least one predetermined event definition; and
 - f) generating a signal if the data meets one of the at least one event definitions.
- 24. The method of claim 19 wherein:

one of the network components is programmed by software and an information technology agent communicates with the software to collect the security data.

20 25. The method of claim 19 wherein:

the step of collecting occurs in real time rather than in batches.

- 26. The method of claim 19 wherein:
- at least two of the plurality of different parser scripts correspond to the same network component.

Serial No. Unknown Applicant: Kathy Maida-Smith, et al. Express Mail No. EL455184579US Atty. Dkt. No. 66997-0102

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NETWORK SECURITY DATA MANAGEMENT SYSTEM AND METHOD

ABSTRACT OF THE DISCLOSURE

A system for compiling security data from an information network includes at least two network components (30,34), each providing data. A data parser (48,52) is coupled to the network components (30,34). The data parser (48,52) has access to two parser scripts that correspond to the 5network components' data. Categorized data can be produced by applying the parser scripts to the data received from the network components (30,34).

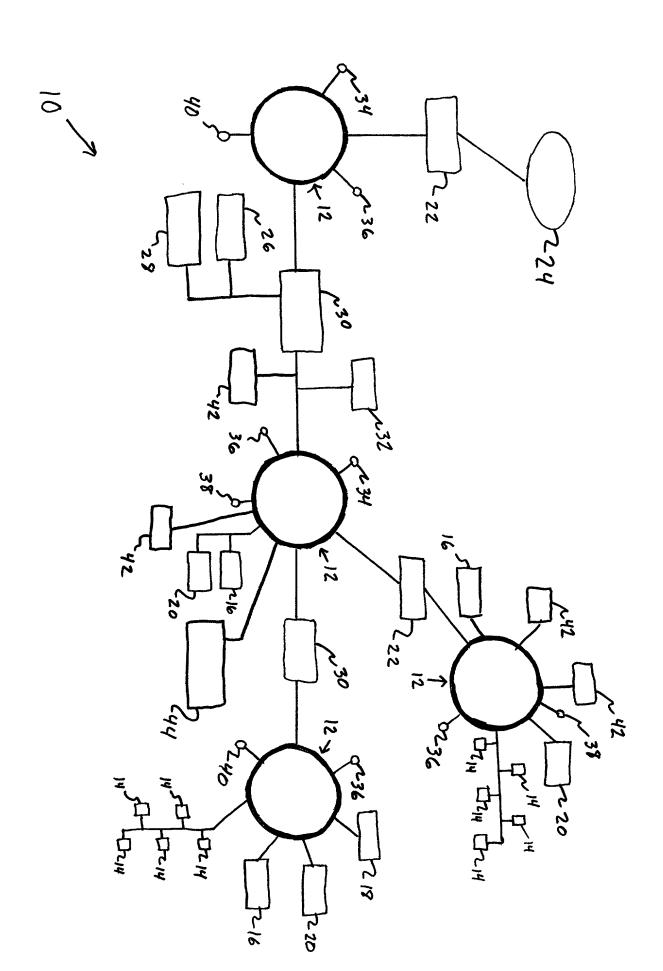
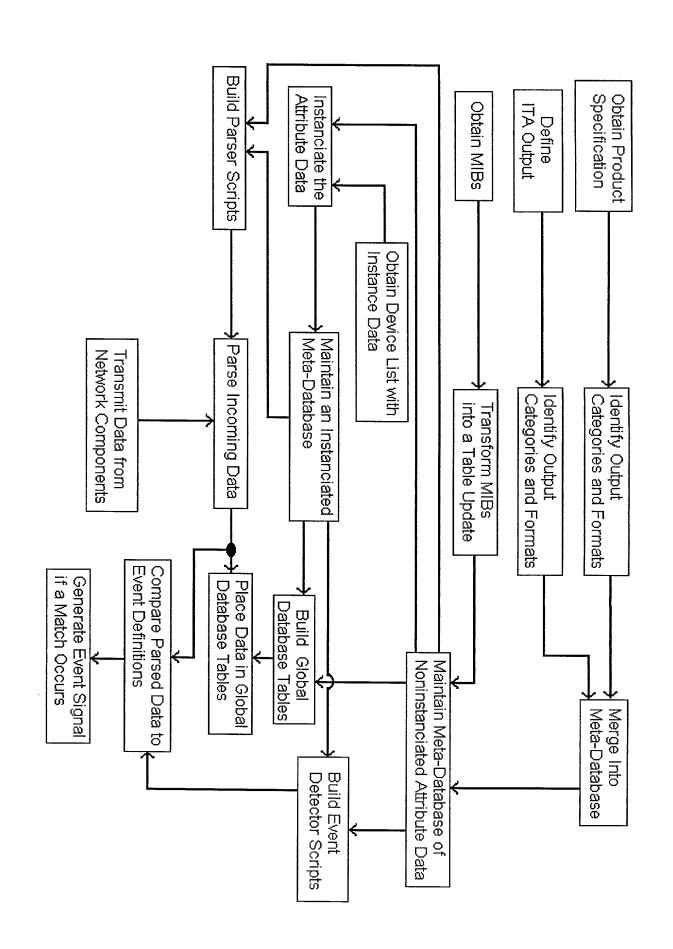


Figure 2



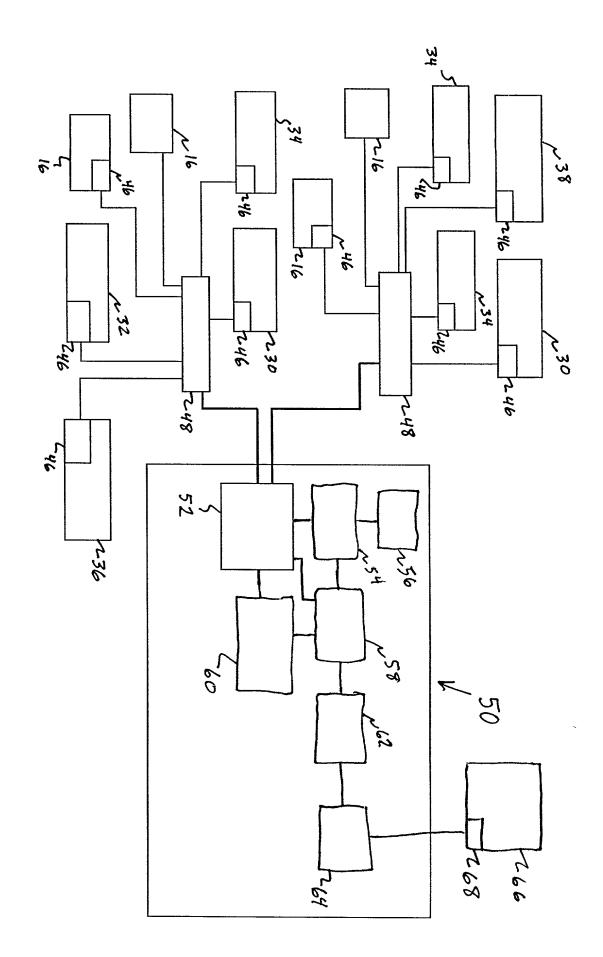


Figure 3

Figure 4 Entrice

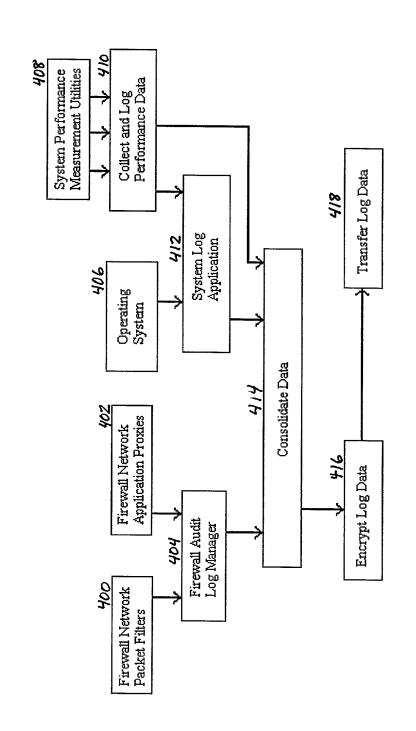


Figure 5

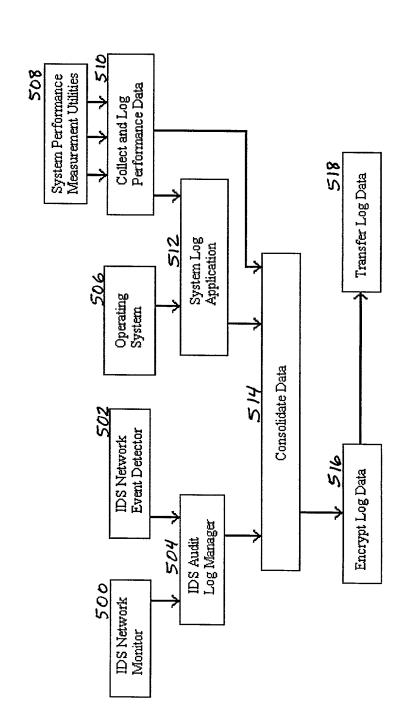


Figure 6

Transaction T	Table 700					
xaction id	name	description		stream id	system type id	1 conditional
X200154	Network Transaction	A record of a network transaction through a CyberGuar	Suard firewall	R100204	T100001	(*Action == "T")
X317189	CPU Utilization	A record of CPU percent idle, percent system, percent I/O, and percent user	ent I/O, and percent user	R113203	T100001	TRUE

	load version revision	CyberGuard 2 2.3
	title	CyberGuard Firewall For UNIX
702	lepoul	d Corp 3412-02
System Type Table	system type id make	T100001 CyberGuar

704 Stream Table

stream id	stream type	stream path	and of record	and of field
R100204	DELIMITED_FLAT_FILE	/var/audit_logs/old/NetguardD	NEWLINE	SPACE
R113203	FIXED_FORM_FLAT_FILE	/store/compressed/stat.log	NEWLINE	4 P

Figure 7

	aggregate	۸_N	۲	۲ 2	۷	۷ Z	۲	۷ V	۷ Z	۷ 2	۲ 2	۲ Z	۸	۲ Z	Υ	AVERAGE	AVERAGE	AVERAGE	AVERAGE	AVERAGE	AVERAGE
	sample period period unit	SYNC N A	SYNC N A	SYNC N A	SYNC N A	SYNC N_A	SYNC N_A	SYNC N_A	SYNC N_A	SYNC N_A	SYNC N_A	SYNC N_A	SYNCIN_A	SYNC/N_A	SYNC N_A	10 MINUTE	10 MINUTE	10 MINUTE	10 MINUTE	10 MINUTE	10 MINUTE
		ASYNC S	ASYNC S	ASYNC S	ASYNC S	ASYNC S	ASYNC S	ASYNC S	ASYNC S	ASYNC S	ASYNC S	ASYNC S'	ASYNC S	ASYNC S'	ASYNC S	1	1	1	1	1	1
	refresh period	/	<i>d</i>	7	<i>y</i>	1	1	<i>y</i>	<i>d</i>	9	,	ď	4	۷ .	ď						
	data unit	۷	۸Ï	۸	۲	۲	۷. V	۲ ۲	۲	۷ 	۸	2 COUNT	COUNT	COUNT	2 COUNT	۲ Z	۲ Z	S PERCENT	S PERCENT	S PERCENT	PERCENT
	legal value	DATE_ISO	TIME_24	CHAR	ALPHANUMERIC	ALPHANUMERIC	1 ₽ d1	15 dl	IP_PROTOCOL	ALPHANUMERIC	ALPHANUMERIC	UNSIGNED_INT_32 COUNT	UNSIGNED_INT_32 COUNT	UNSIGNED_INT_32 COUNT	UNSIGNED_INT_32 COUNT	DATE_ISO	TIME 24	UNSIGNED_INT_16 PERCENT	UNSIGNED_INT_16 PERCENT	UNSIGNED_INT_16 PERCENT	UNSIGNED_INT_16 PERCENT
	description	Date of network transaction	Time of network transaction	Transaction Type	Name of source firewall network interface	Name of destination firewall network interface	Source IP address of packet	Destination IP address of packet	IP protocol of packet	Source port number / service name of packet or ICMP type	Destination port number / service name of packet or ICMP type	Number of packets passed from source to destination	Number of packets passed from destination to source	Number of bytes passed from source to destination	Number of bytes passed from destination to source	Date of measurement	Time of measurement	Percent CPU Idle	Percent of CPU used for system tasks	Percent of CPU used for I/O tasks	Percent of CPU used for user tasks
	name	Date	Time	Action	Source_IF	Destination_IF	Source_Address	Destination_Address	Protocol	Source_Port	Destination_Port	Packets_Sent	Packets_Received	Bytes_Sent	Bytes_Received	Date	Time	CPU Idle	System Percentile	I/O Percentile	User Percentile
le 706	element id	E312073	E312074	E312075	E312076	E312077	E312078	E312079	E312080	E312081	E312082	E312083	E312084	E312085	E312086	E444120	E444121	E444122	E444123	E444124	E444125
Element Table	xaction id	X200154	X200154	X200154	X200154	X200154	X200154	X200154	X200154	X200154	X200154	X200154	X200154	X200154	X200154	X317189	X317189	X317189	X317189	X317189	X317189

Delimited Flat File Table	708						
element id	encoding	record num	field num	subfield delim	subfield num	start in field	end in field
E312073	ASCII_7_BIT	ALL	1	NONE	0	0	0
E312074	ASCII_7_BIT	ALL		2 NONE	0	0	•
E312075	ASCII_7_BIT	ALL	e e	ONE	0	0	0
E312076	ASCII_7_BIT	ALL	1 4		1	0	0
E312077	ASCII_7_BIT	ALL	/ 7		2	0	0
E312078	ASCII_7_BIT	ALL		5 NONE	0	0	0
E312079	ASCII_7_BIT	ALL		ONE	0	0	0
E312080	ASCII 7 BIT	ALL		7 NONE	0	0	0
E312081	ASCII_7_BIT	ALL		ONE	0	0	0
E312082	ASCIL 7_BIT	ALL		9 NONE	0	0	0
E312083	ASCII_7_BIT	ALL		ONE	0	0	0
E312084	ASCII_7_BIT	ALL		ONE	0	0	0
E312085	ASCII_7_BIT	ALL	12 N	12 NONE	0	0	0
E312086	ASCII_7_BIT	ALL	13 N	ONE	0	0	ō

	encoding	record num	start in record	end in record
E444120	ASCII_7_BIT	ALL		J.
E444121	ASCII_7_BIT	ALL	12	¥
E444122	ASCII_7_BIT	ALL	. 21	2.
E444123	ASCII_7_BIT	ALL	24	2(
E444124	ASCII 7 BIT	ALL	. 27	7
E444125	ASCIL_7_BIT	ALL	30	æ.

Estable Figure 8

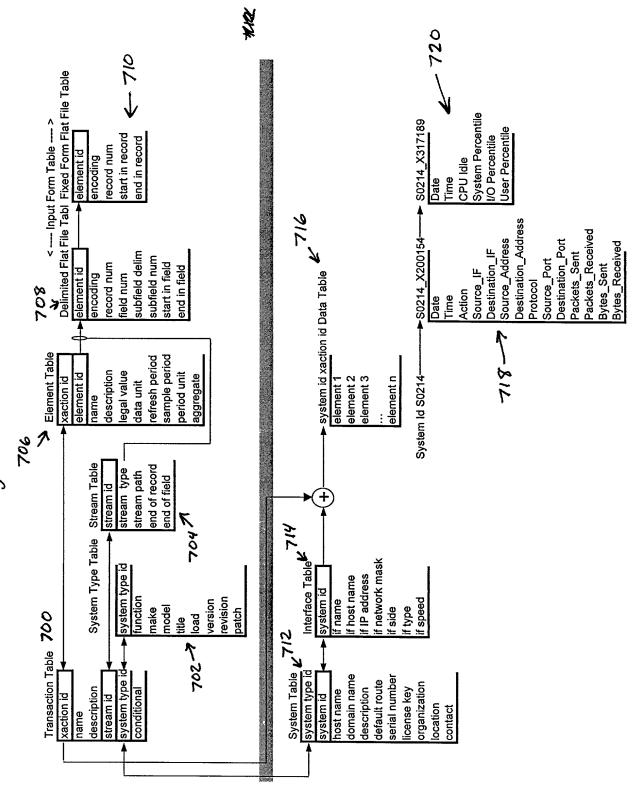


Figure 9

1999/10/2709:00:00043004010023 1999/10/2709:10:00022003005010 1999/10/2709:20:00043004010023 1999/10/2709:30:00043004010023 1999/10/2709:40:00043004010023 1999/10/2709:50:00043004010023

Figure 10

S0214_X317189 Data Table

Date	Time	CPU Idle	System Percentile	I/O Percentile	User Percentile
1999/10/27	09:00:00	43	4	10	23
1999/10/27	09:10:00	22	3	5	10
1999/10/27	09:20:00	43	4	10	23
1999/10/27	09:30:00	43	4	10	23
1999/10/27	09:40:00	43	4	10	23
1999/10/27	09:50:00	43	4	10	23

Attorney Docket No.: 66997-0102

DECLARATION AND POWER OF ATTORNEY FOR PATENT APPLICATION

As the below named inventor, I hereby declare that:

My residence, post office address and citizenship is as stated below next to my name;

		Network Security Data		, the
specification of which				
\square is attached heret	ю.			
was filed on No	vember 22, 1999 as	Application Number	and	was
amended on	(if applicable)	·		
amended by any amend to patentability as defi	dment referred to aboned in 37 C.F.R. §1.5		luty to disclose inform	icluding the claims, a nation which is materia
foreign application(s) f designated at least on	ign priority benefits uor patent or inventor's e country other than plication(s) for pater	rior Foreign Application(under Title 35, United S s certificate, or § 365 (a the United States of Ar nt or inventor's certific	States Code, §119(a)- a) of any PCT internat merica, listed below a	ional application whic Ind have also identifie
Country	Application Number	Date of Filing (day, month, year)	Date of Issue (day, month, year)	Priority Claimed
200				
				Yes No D
				Yes No

Prior Provisional Application(s)

I hereby claim the benefit under Title 35, United States Code § 119(e) of any United States provisional application(s) listed below:

Application Number	Date of Filing (day, month, year)

Prior United States Application(s)

I hereby claim the benefit under Title 35, United States Code, §120 of any United States application(s), or § 365(c) of any PCT international application designating the United States of America, listed below and, insofar as the subject matter of each of the claims of this application is not disclosed in the prior United States application in the manner provided by the first paragraph of Title 35, United States Code, §112, I acknowledge the duty to disclose material information as defined in Title 37, Code of Federal Regulations, §1.56(a) which occurred between the filing date of the prior application and the national or PCT international filing date of this application:

Application	Date of Filing	Status - Patented,
Number	(day, month, year)	Pending, Abandoned

And I hereby appoint, both jointly and severally, as my attorneys with full power of substitution and revocation, to prosecute this application and to transact all business in the Patent and Trademark Office connected herewith the following attorneys, their registration numbers being listed after their names:

Michael Hawes, Registration No. 38,487, Rodger L. Tate, Registration No. 27,399; Scott F. Partridge, Registration No. 28,142; Jerry W. Mills, Registration No. 23,005; Claude E. Cooke, Jr., Registration No. 34,142; Mitch D Lukin, Registration No. 30,772; Roger Fulghum, Registration No. 39,678; James Remenick, Registration No. 36,902; Lori D. Stiffler, Registration No. 36,939; Melissa Szanto, Registration No. 40,834; Michael Wilson, Registration No. 32,835; and Alan Witte, Registration No. 36,061.

All correspondence and telephone communications should be addressed to: Baker & Botts, L.L.P., The Warner, One Shell Plaza, 910 Louisiana, Houston, Texas 77002, telephone number (713) 229-1234, which is also the address and telephone number of each of the above listed attorneys.

I hereby declare that all statements made herein of my own knowledge are true and that all statements made on information and belief are believed to be true; and further that these statements were made with the knowledge that willful false statements and the like so made are punishable by fine and imprisonment, or both, under 18 U.S.C. § 1001, and that such willful false statements may jeopardize the validity of the application or any patent issuing thereon.

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